

SPLIT-NEIGHBOURHOOD OF A GRAPH

Shashidhar G. Kumbar

Assistant Professor

Department of Mathematics

Govt. First Grade College, Muddebihal

Karnataka

ABSTRACT

A Neighbourhood set S of a graph G is a Split Neighbourhood set if the induced subgraph $\langle V - S \rangle$ is disconnected. The split-Neighbourhood number $n_s(G)$ is the minimum cardinality of a split-Neighbourhood set. In this paper, we have obtained bounds for $n_s(G)$ in terms of order, size and other parameters of graphs.

Keywords: Domination number, Split Domination Number, Split-Neighbourhood.

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1. INTRODUCTION

The graphs considered here are finite, undirected, without loops or multiple edges and connected. Unless otherwise stated, all graphs are assumed to have ' p ' vertices and ' q ' edges.

A set S of vertices in graph G is a Neighbourhood set (n -set) of G if $\bigcup_{u \in S} N(u) = V$.

$\langle N(u) \rangle$, where $\langle N(u) \rangle$ is the subgraph induced by u and all vertices adjacent to $u \in S$, $\{u\}$ is not Neighbourhood set of G . The Neighbourhood number $n(G)$ of G is a

minimum cardinality of a n – set of G . This parameter is introduced by E. Sampathkumar and P. S. Neeralagi [6].

There are many types of domination numbers in literature [2]. Similarly we can define different types of Neighbourhood numbers by imposing certain conditions on Neighbourhood sets and derive some of the properties.

A Neighbourhood set S is said to be a maximal Neighbourhood set of G if the induced subgraph $\langle V - S \rangle$ is not a Neighbourhood set of G . The maximal Neighbourhood number $n_m(G)$ of G is the minimum cardinality of a maximal Neighbourhood set of G . This parameter is introduced by N.D. Soner et al [6].

In this chapter, we introduce the concept of SplitNeighbourhood as follows :

A Neighbourhood set S of a graph G is a Split Neighbourhood set if the induced subgraph $\langle V - S \rangle$ is disconnected. The SplitNeighbourhood number $n_s(G)$ is the minimum cardinality of a SplitNeighbourhood set.

Thus, we observe that for any graph G ,

$$(G) \leq n^\circ(G) \leq n_s(G) \leq \alpha^\circ(G) \dots\dots\dots (I)$$

$$(G) \leq \gamma_s(G) \leq n_s(G) \leq \alpha^\circ(G) \dots\dots\dots (II)$$

Now we will prove the following results.

2. RESULTS

Theorem A [4] A dominating set D of G is a Split dominating set if and only if there exists two vertices $w_1, w_2 \in V - D$ such that $w_1 - w_2$ path contains a vertex of D .

Theorem 2.1 For any graph G , $n^\circ(G) \leq n_s(G) \dots\dots\dots (1)$

Further the bound is attained if and only if there exists two vertices $w_1, w_2 \in V - S$ such that every $w_1 - w_2$ path contains a vertex of S where S is a n° – set of G .

Proof: Equation (1) follows from the definition of SplitNeighbourhood set.

Further let S be a Neighbourhood set such that there exists two vertices $w_1, w_2 \in V - S$ such that every $w_1 - w_2$ path contains a vertex of S . Then $\langle V - S \rangle$ is disconnected. Hence S is a SplitNeighbourhood set. This implies $n_s(G) \leq n^\circ(G)$. Then from (1) we have $n^\circ(G) = n_s(G)$.

Conversely suppose the bound is attained. Then if S is a Neighbourhood set, it is also a SplitNeighbourhood set. This implies $\langle V - S \rangle$ is disconnected. Hence there exist two vertices $w_1, w_2 \in V - S$ such that every $w_1 - w_2$ path contains a vertex of S .

Theorem B [6] :For a graph G , $n^\circ(G) = \gamma(G)$ if and only if there exists a minimum dominating set S . Such that every line in $\langle V - S \rangle$ belongs to $\langle (u) \rangle$ for some $u \in D$.

Theorem 2.2 For any graph G ,

$$\gamma_s(G) \leq n_s(G) \dots \dots \dots (2)$$

Further the bound is attained if and only if there exists a minimum Split dominating set S such that every line in $\langle V - S \rangle$ belongs to $\langle N(u) \rangle$ for some $u \in S$.

Proof :Since every SplitNeighbourhood set is a Split dominating set, hence Split dominating number is less than SplitNeighbourhood number. Suppose the bound is attained. This implies the condition is satisfied from Theorem 5.A [4].

Conversely, suppose that given condition is satisfied for some Split dominating set S . Then again by Theorem 5.B [6], S is a Neighbourhood set. Since $\langle V - S \rangle$ is disconnected. S is a Split Neighbourhood set and hence from (2) the bound is attained.

Theorem C [6] For any graph G without isolated points,

$$\gamma(G) \leq n^\circ(G) \leq \alpha^\circ(G)$$

Theorem 2.3 For any graph G without isolated points,

$$n_s(G) \leq \alpha^o(G) \dots\dots\dots (3)$$

Further the bound is attained if and only if there exist a Split Neighbourhood set S of G for which $V - S$ is independent with at least two vertices.

Proof : Let S be vertex cover of G . Then, $V - S$ is independent with at least two vertices. This implies, $\langle V - S \rangle$ is disconnected. Also S is a Neighbourhood set from Theorem 5.C [6]. Hence S is a Split Neighbourhood set of G . This proves that the Split Neighbourhood number is less than or equal to vertex covering number.

Now to prove the second part, suppose there exist a Split Neighbourhood set S of G for which $V - S$ is independent with at least two vertices. This implies S is a vertex cover of G . Thus vertex covering number of G is less than or equal to the cardinality of S . Hence from (3), the bound is attained.

Conversely, suppose equality holds. Then there exists a Split Neighbourhood set S which is a vertex cover with $|S| = \alpha^o(G)$. Then obviously $V - S$ is independent with at least two vertices.

Theorem D [4] For any graph G , $\gamma \leq \gamma_s$

Hence from Theorem 5.1, 5.2, 5.3, 5.C [6] and 5.D [4]

we have,

$$(G) \leq n^o(G) \leq n_s(G) \leq \alpha^o(G) \dots\dots\dots (I)$$

$$(G) \leq \gamma_s(G) \leq n_s(G) \leq \alpha^o(G) \dots\dots\dots (II)$$

Theorem 2.4 For any graph G ,

$$(G) \leq n_s(G) \dots\dots\dots (4)$$

Where (G) is the connectivity of graph G .

Proof : Let S be a Split Neighbourhood set of G . Then $\langle V - S \rangle$ is disconnected.

Hence $(G) \leq n_s(G)$

Next, we list the exact value of $n_s(G)$ for some standard graphs

Theorem 2.5 (i) For a path P_n with n vertices,

$$n_s(P_n) = \left\lfloor \frac{n}{3} \right\rfloor \quad n \geq 3 \dots \dots \dots (5)$$

(ii) For a circle C_n with n vertices,

$$n_s(C_n) = \left\lfloor \frac{n}{2} \right\rfloor \quad n \geq 4 \dots \dots \dots (6)$$

(iii) For a wheel W_n with n vertices,

$$n_s(W_n) = 3 \quad n \geq 5 \dots \dots \dots (7)$$

(iv) For a bipartite graph, without isolates, with bipartition $\{v_1, v_2\}$

of $V(G)$,

$$n_s(G) \leq \min \{|v_1|, |v_2|\} \dots \dots \dots (8)$$

Moreover the bound is attained by the graphs $K_{m,n}$

Proof :

(i) For a path P_n with n vertices where $n \geq 3$, every Neighbourhood set is a Split Neighbourhood set. Hence (5) follows.

(ii) For a cycle C_n with n vertices where $n \geq 4$, every Neighbourhood set is a Split Neighbourhood set. Hence (6) follows.

(iii) For a wheel W_n with n vertices where $n \geq 5$, the vertex with degree $p - 1$ together with two non adjacent vertices on the cycle form a Split Neighbourhood set. Hence (7) follows.

(iv) For a bipartite graph with bipartition $\{V_1, V_2\}$ of $V(G)$, both the sets with cardinality V_1 and V_2 are Split Neighbourhood sets. Hence (8) follows. Further if it is a complete bipartite graph then equality holds since for any V_i , $i=1,2,3,\dots\dots$

$V_i - \{u\}$ is not a Split Neighbourhood set.

Theorem E [6] For any bipartite graph G without isolated points,

$$n^\circ(G) = \alpha^\circ(G) = \beta_1(G)$$

Theorem 2.6 For any bipartite graph G without isolated points,

$$n^\circ(G) = n_s(G) = \alpha^\circ(G) = \beta_1(G) \dots\dots\dots (9)$$

Proof : This follows from Theorem 5.E [6] and Result (I)

Theorem 2.7 A Split Neighbourhood set S is minimal if and only if for each vertex $v \in S$, one of the following conditions is satisfied

- (i) v is an isolate in $\langle S \rangle$
- (ii) There exist a vertex $u \in V - S$ adjacent to v but not adjacent to any vertex $w \in S$ adjacent to v .
- (iii) $\langle (V - S) \cup \{v\} \rangle$ is connected.

Proof: Suppose S is minimal, on the contrary, if there exists $v \in S$ such that v does not satisfy any of the given conditions. Then $S' = S - \{v\}$ is a Neighbourhood set of G from (i) and (ii) and $\langle V - S' \rangle$ is disconnected from (iii) This implies S' is Split Neighbourhood set of G . This is a contradiction. This proves that necessity.

Sufficiency is straight forward.

Theorem F [1] : For any non trivial connected graph G ,

$$\alpha^\circ(G) + \beta^\circ(G) = p$$

Theorem 2.8:

i) For any graph G ,

$$(G) \leq n^\circ(G) \leq n_s(G) \leq (\chi(G) - 1)\beta^\circ(G) \dots\dots\dots (10)$$

Provided $(G) \geq 2$, where $\chi(G)$ is the chromatic number of graph G .

ii) If G is bipartite graph which is not totally disconnected, Then,

$$(G) \leq n^{\circ}(G) \leq n_s(G) \leq \beta^{\circ}(G) \leq \chi(\bar{G}) \dots\dots\dots (11)$$

Where $\bar{G}$ is complement of G .

Proof : Here we need to establish only the upper bound since lower bounds from I.

From Theorem 5.F [1] and the fact that $p \leq (G)(\beta^{\circ}(G))$

(See [1]) we have,

$$p - \beta^{\circ}(G) \leq \beta^{\circ}(G)(\chi(G) - 1)$$

$$\text{i.e. } \alpha^{\circ}(G) \leq \beta^{\circ}(G)(\chi(G) - 1)$$

Hence (10) follows from (1) and the fact that $\alpha^{\circ}(G) \leq \beta^{\circ}(G)(\chi(G) - 1)$

If G is bipartite, $(G) = 2$. Also (10) implies $n_s(G) \leq \beta^{\circ}(G)$

Hence (11) follows from the facts that $n_s(G) \leq \beta^{\circ}(G)$ and $\beta^{\circ}(G) \leq \chi(\bar{G})$ (See [1]).

Theorem 2.9 For any graph G ,

$$n_s(G) = 1 \dots\dots\dots (12)$$

If and only if there exists a cut vertex with degree $p - 1$

Proof : Suppose v is cutvertex of G of degree $p - 1$, then $\{v\}$ is a Neighbourhood set. Further since $\langle V - \{v\} \rangle$ is disconnected. This implies $\{v\}$ is a Split Neighbourhood set. Hence $n_s(G) = 1$

Conversely, suppose $n_s(G) = 1$. Then, obviously there exists a cutvertex which is adjacent to all vertices. Hence there exists a cutvertex with degree $p - 1$.

Theorem G [6] For any (p, q) graph G ,

$$p - q + q^{\circ} \leq n^{\circ}(G) \leq p - \Delta(G)$$

$$\lceil \frac{p}{\Delta(G)+1} \rceil \leq n^{\circ}(G) \leq p - \beta^{\circ}(G) + p^{\circ}$$

Where q° = minimum $\{q(\langle D \rangle); D \text{ is a minimal dominating set of } G\}$

p° = the number of isolated vertices in G ,

β° = set of independent vertices in G .

Theorem 2.10 For any connected (p, q) graph G ,

$$p - q + q^{\circ} \leq n_s(G) \dots\dots\dots (13)$$

$$\lceil \frac{p}{\Delta(G)+1} \rceil \leq n_s(G) \leq p - \beta^{\circ}(G) \dots\dots\dots (14)$$

Proof : The lower bounds in (13) and (14) follow from (1) and Theorem 5.G [6]. To prove upper bound in (14), we observe that $(V - M)$ is a Split Neighbourhood set where M is the set of β° independent points of G .

The lower bound in (13) and (14) is attained for the following graph in Figure 5

The upper bound in (14) is attained for any tree

The lower bound in (14) is attained by the following graph in figure 6.

Theorem 2.11

(i) $n_s(G) > p - \Delta(G)$ if there exist a non-cutvertex of degree $p - 1$

(ii) $n_s(G) \leq p - \Delta(G)$ if G has no triangle.

Proof :

(i) Let G has a non-cutvertex v of degree $p - 1$. Then $\Delta(G) = p - 1$. Since v is the non-cutvertex, $n_s(G) \geq 2$. Hence $n_s(G) > p - \Delta(G)$.

(ii) If G has no triangle then $n_s(G) \leq p - \Delta(G)$ from (9) and Theorem 5.G [6].

Now we obtain a Nordhaus-Gaddum type result.

Theorem 2.12 Let G be a graph such that both G and $\bar{G}$ are connected, then

$$n_s(G) + n_s(\bar{G}) \leq p(p-3) \dots\dots\dots (15)$$

Further the bound is attained if and only if $G = P_4$

Proof : We have $n_s(G) \leq \alpha^\circ(G)$ from (3).

Since both G and $\bar{G}$ are connected, $\Delta(G), \Delta(\bar{G}) < p-1$

This implies $\beta^\circ(G), \beta^\circ(\bar{G}) \geq 2$.

Hence $n_s(G) \leq p-2$

$$= 2(p-1) - p$$

$$\leq (2q-p)$$

Similarly $n_s(\bar{G}) \leq 2\bar{q} - p$

Thus $n_s(G) + n_s(\bar{G}) \leq 2(q + \bar{q}) - 2p$

$$\leq (p-1) - 2p$$

$$= (p-3)$$

Suppose the bound is attained, then $n_s(G) = 2q - p$ and $n_s(\bar{G}) = 2\bar{q} - p$. This implies q and $\bar{q} < p$. Hence G and $\bar{G}$ are trees. i.e. $G = P_4$

Now we will establish a relation between Split Neighbourhood number and maximum Neighbourhood number.

Theorem 2.13 Let G be a graph with $\beta^\circ(G) \geq 3$ and possess no triangles.

Then, $n_s(G) \leq n_m(G) \dots\dots\dots (16)$

Proof : Let S be a maximal Neighbourhood set of G . Then $\langle V - S \rangle$ is totally disconnected with at least two vertices. Thus S is a Split Neighbourhood set. Hence (16) holds.

Theorem H [7] For any graph G ,

$$n_m(G) \leq \alpha^o(G) + 1$$

Theorem 2.14 Let G be a graph without triangle, then

$$n_m(G) \leq n_s(G) + 1 \dots\dots\dots(17)$$

Proof : The Proof of (17) follows from (9) and Theorem 5.H [7].

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